

## Scoreboard Method

Functional Unit Status Table [Add  $R_2, R_3, R_4$ ]

Unit ID	Unit Name	Busy	Destination Register $R_d$	Source Registers			
				$R_{s1}$	Ready	$R_{s2}$	Ready
1	Load/Store	NO					
2	Multiplier	YES	R5	R1	YES	R2	YES
3	Adder_1	NO					
4	Adder_2	YES	R2	R2	YES	R5	NO

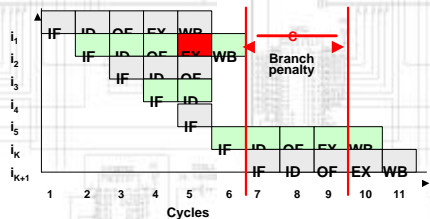
## Scoreboard Method

Destination Register Status Table [Add  $R_2, R_3, R_4$ ]

	R1	R2	R3	R4	R5	R6
Unit Id		4			2	

## Control Hazards

⊕ All branch instructions have the potential to drain the pipe, imposing a branch penalty  $c$  which is proportional to the overall pipelength



## Branching

- ⊕ Initially, of course, we cannot anticipate branches
- ⊕ Once encountered we can preserve useful information:
  - ⊕ Branch location in memory (PC)
  - ⊕ Branch history (ie. taken or not)
  - ⊕ Branch category

## Branch Category

- ⊕ Unconditional (JUMP, GOSUB)
- ⊕ Loop branch
  - ⊕ If we branch back not far from the branch location we probably have a loop body
  - ⊕ Preserve it in the IF stage
- ⊕ Conditional branch
  - ⊕ Use history to decide between target or no branch

## Branch Penalty

Average number of cycles per inst.

with  $p_b$  = probability instruction is a branch

$$t_{ave} = p_b(\text{average number of cycles per branch instruction}) + (1 - p_b)(\text{average number of cycles per non-branch inst.})$$

Average number of cycles per branch instruction

If target path =  $1 + c$  cycles with  $c$  as branch penalty

If sequential = 1 cycles

With  $p_t$  = probability branch is taken

Average number of cycles per branch instruction =

$$p_t(1+c) + (1 - p_t)1$$

$$t_{ave} = p_b(p_t(1+c) + (1 - p_t)1) + (1 - p_b)(1)$$

$$= p_b p_t c + 1$$

## Branch Penalty

$$t_{ave} = p_b p_t c + 1$$

Typically  $0.1 < p_b < 0.3$  and  $0.6 < p_t < 0.7$

For a 5-stage pipe with  $c=3$

$$t_{ave} = p_b p_t c + 1 = 0.2 * 0.65 * 3 + 1 = 1.39$$

$$\text{Efficiency} = 1 / t_{ave} * 100 = 72\%$$

## Impact of Branches

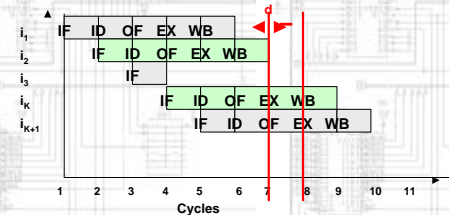
- ⊕ Techniques for reducing the impact of branches
  - ⊕ Branch Prediction
  - ⊕ Multiple Prefetching
  - ⊕ Delayed Branching

## Branch Prediction

- ✦ Predict the outcome of the branch before it is executed
- ✦ Incorrect predictions may increase penalties
- ✦ Static prediction -
  - ✦ for example, always assume a branch will be taken
- ✦ Dynamic -
  - ✦ use branch target buffer (BTB) (see 3BA4)

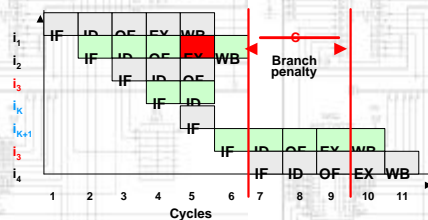
## Control Hazards

Penalty for a **correctly** chosen target path



## Control Hazards

Penalty for a **incorrectly** chosen target path



## Performance effect of Branch Prediction

with  $p_t$  = probability branch is taken

Average number of cycles per branch instruction =  $p_t(1+c) + (1-p_t)1$

with  $p_r$  = probability of right prediction ->

Average number of cycles per branch instruction =

$$p_r \underbrace{[p_t(1+d) + (1-p_t)1]}_{\text{Correct Prediction}} + (1-p_r) \underbrace{[p_t(1+c) + (1-p_t)(1+c)]}_{\text{Incorrect Prediction}}$$